

Natural Information Systems

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Natural

- An object/arrow is defined uniquely over 3 levels up to natural isomorphism
- 3 levels of arrows:
 - Object A in Category C , $1_A : A \rightarrow A$
 - Functor $F: C \rightarrow C$; $G: C \rightarrow C$
 - Natural Transformation $\alpha: F \rightarrow G$
- Information System:
 - As in Universe

Current Information Systems

Existing approach is models based on a standard mathematical structure:

Hierarchical

Trees

Network

Graphs

Relational

Sets

Functional

Functions

Object-oriented

Objects/methods

What about Categories?

- Previous work:
 - Rosebrugh and co-workers
 - Spivak
 - Baclawski
- Has largely concentrated on:
 - Representing relational databases or the entity-relationship model in categories
 - Sketches popular
 - Developed by Charles Ehresmann

Sketch outline

- A sketch is a
 - Graph with
 - a set of diagrams
 - a set of cones defining which diagrams have limits
 - a set of cocones defining which diagrams have colimits
- Compared to Cartesian closed, a sketch relaxes
 - The terminal object requirement
 - The need for all diagrams to have limits and colimits
- But a sketch is limited to graphs

Categorification

- Use of sketches is categorification
 - Transforming existing techniques on a 1:1 basis from application to categories
 - May provide useful support for current database model research.
- That does not realise the full potential for category theory in advancing database techniques

Fundamental Approach

- Start from basics
 - a clean sheet
- Decide on requirements for an information system
- Identify features of category theory that help to meet requirements
- Produce a framework which satisfies software engineering principles:
 - High cohesion
 - Low coupling

Candidate - Topos

- Categorical structure
- Attracts much attention in standard texts
 - Mac Lane (CWM), Goldblatt (Topoi), Johnstone (Topos Theory)
- Captures properties of sets
- Based on Cartesian closed categories (CCC)
 - Basis of much of our recent work
 - Also dealt with fully in standard texts

The Topos - Definition

- A category \mathcal{C} that
 - is finitely complete
 - Limits of all finite diagrams (cones)
 - is finitely co-complete
 - Colimits of all finite diagrams (cocones)
 - Follows automatically if finitely complete and CCC
 - is Cartesian closed
 - has a subobject classifier

Example of CCC

- Pullback

- Terminal object 1

- Exactly one arrow from every object in category to the terminal object (least upper bound)

- Products

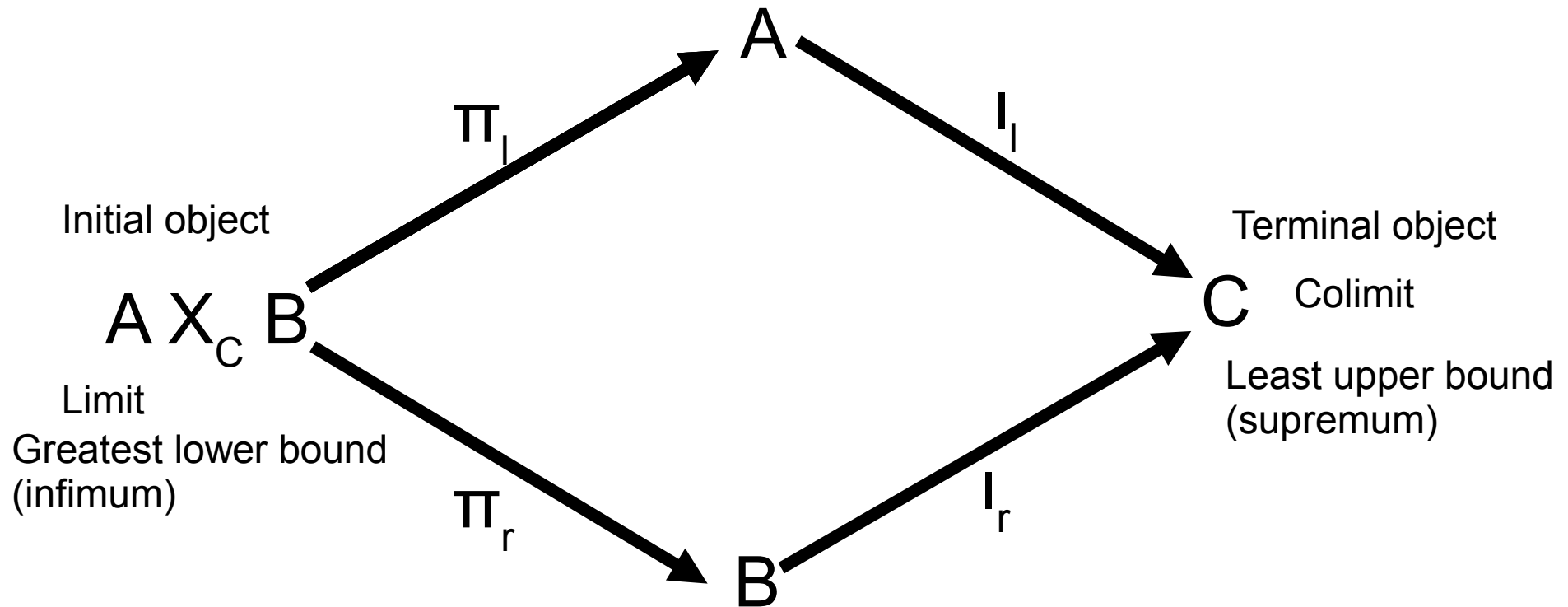
- All objects A,B are related through products $A \times B$

- Exponentiation (connectivity)

- $\text{hom}(A \times B, C) \equiv \text{hom}(A, C^B)$

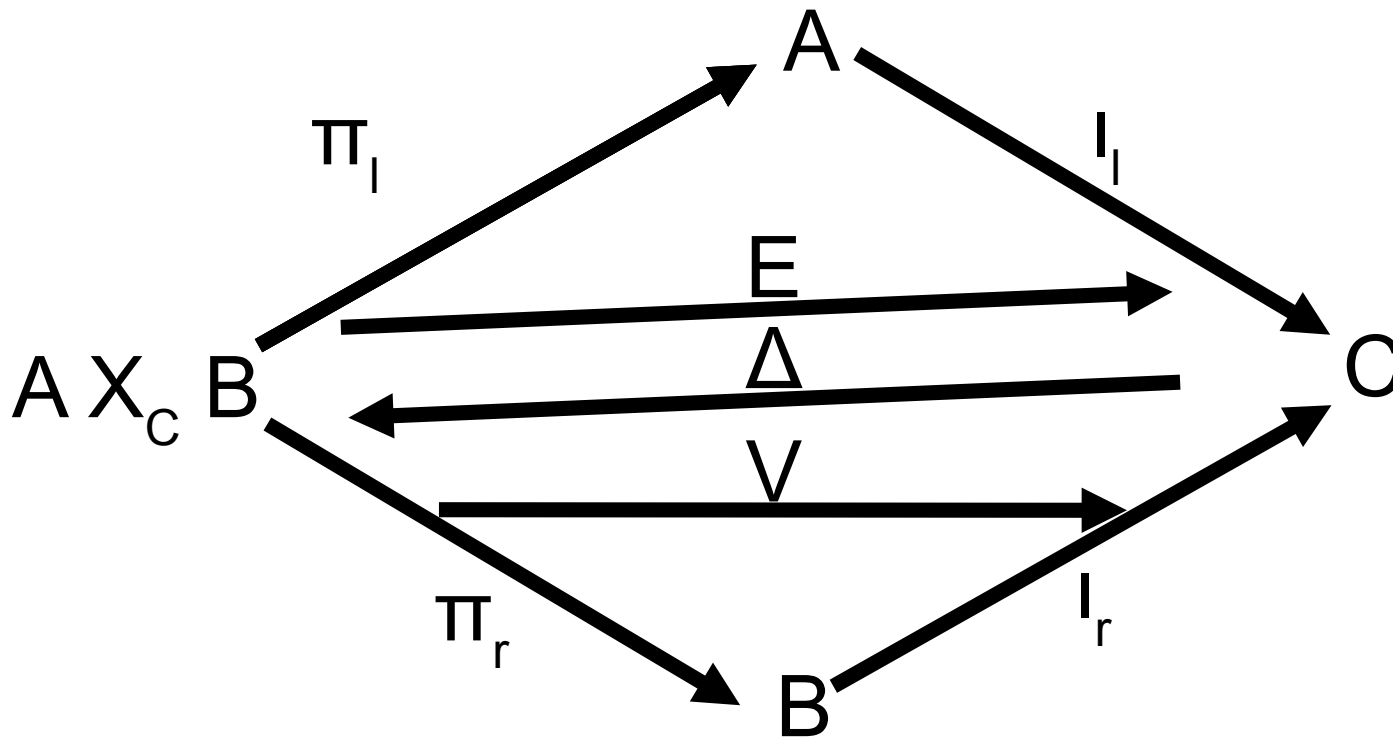
- $F: _ \times B: \zeta \dashrightarrow \zeta; \quad G: _ ^B: \zeta \rightarrow \zeta; \quad F \dashv G$

Pullback – Cartesian closed Category



C is $A+B+C$

Pullback Logic



Adjointness requirements $E \dashv \Delta$ and $\Delta \dashv V$

Build-up of Logic: E existential quantifier, V universal quantifier, Δ diagonal functor, Π projection; i inclusion

Arguments for Topos

- The natural categorial structure for information systems is the topos with its:
 - cartesian product for representing relationships
 - terminal object for identity
 - subobject classifier for membership criteria
 - internal logic Heyting for query and rule processing
 - internal logic is object oriented
 - Heyting is intuitionistic, more general than Boolean

Arguments against Topos

- No readily-accessible examples of usage (applications)

Subobject Classifier

- Live within the topos
- Defined by pullback square:

$$\begin{array}{ccc} U & \longrightarrow & 1 \\ \downarrow j & & \downarrow \text{true} \\ X & \xrightarrow{X_j} & \Omega \end{array}$$

- Where Ω is the subobject classifier
 - 1 is the terminal object of the topos
 - $j: U \rightarrow X$ is an arrow in \mathcal{C}
 - X_j is the characteristic function
 - U is the limit of the pullback, X is the subobject

Subobject Classifier Example 1

- Defined by commuting pullback square:

$$\begin{array}{ccc} U & \longrightarrow & 1 \{t\} \\ \downarrow j & & \downarrow \text{true} \\ S & \xrightarrow{X_j} & \Omega \{0, 1\} \end{array}$$

Subobject classifier is Boolean $\{0, 1\}$

Characteristic function X_j defines subobject S of category represented by 1

Subobject Classifier Example 2

- Defined by commuting pullback square:

$$\begin{array}{ccc} U & \longrightarrow & 1 \\ \downarrow j & & \downarrow \text{true} \\ S(\Omega \times \Omega) & \xrightarrow{X_j} & \Omega \text{ (power-object)} \end{array}$$

Subobject classifier is non-Boolean, a power-object of some objects

Characteristic function X_j defines subobject S of category represented by 1

S is of type AND (intersection)

Topos Logic

- Quantification, projection, product, join through pullback diagrams
- Mitchell-Bénabou Language of a Topos
 - Types are defined + variables of the types
 - Formulae are defined to build expressions
 - Predicates are constructed for membership tests
 - Logical operations include: intersection, union,
 - Internal logic is intuitionistic (Heyting)
 - Handling of negation is more sophisticated
 - e.g. 'not unhappy'

Requirements/Capabilities

Approach	Maths structure	Structuring Capability	Searching	Query symmetry	Query Closure	Transactions	Interoperability	Commercial Examples
Hierarchical	Trees	1:N	Tree traversal	No, bias to downwards direction	No, tabular display	Yes, CICS	No	IMS
Network	Graphs	N:M	Graph traversal	No, bias to initially defined paths	No, tabular display	Yes	No	IDMS
Relational	Relations of sets	N:M	Set operations	Yes	Yes	Yes	No	Oracle, DB2, Access
Functional	Functions	1:N	Function composition	No, bias to initially defined paths	No	No	No	None
Cartesian-closed	Pullbacks	N:M	Quantifiers	Yes	No	Yes, monad	Yes, natural	None
Topos	Categorical Topos	N:M	Quantifiers, internal Heyting logic	Yes	Yes	Yes, monad	Yes, natural	None

Visualisation Application

Vickers, Paul, Faith, Joe, & Rossiter, Nick,

Understanding Visualization: A Formal Approach using
Category Theory and Semiotics,

IEEE Transactions On Visualization And Computer
Graphics 158. IEEE Transactions On Visualization And
Computer Graphics, 2012 Jun;19(6):1048-61. doi:
10.1109/TVCG.2012.294 (2013). pdf

Initial Category for Visualisation

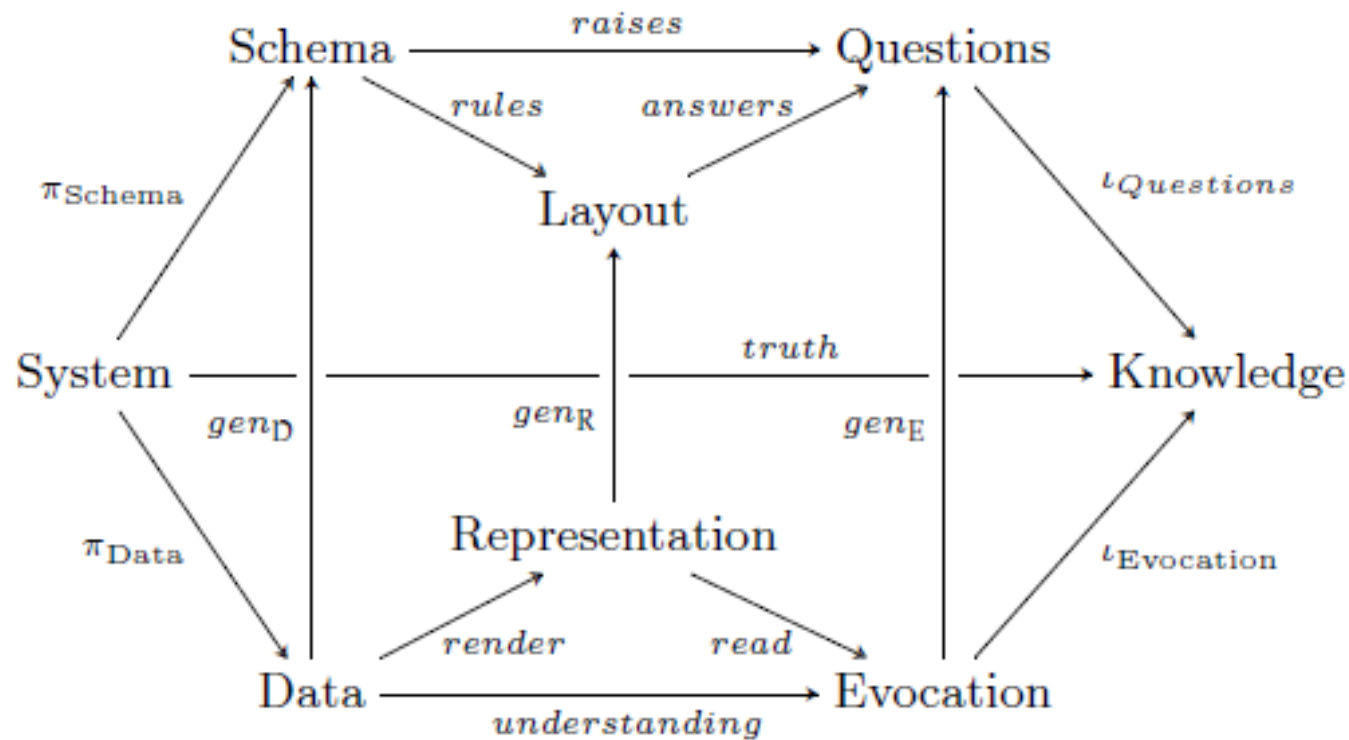


Fig. 3. The visualization process category with product System and co-product Knowledge. System has projection morphisms to Schema and Data. Questions and Knowledge have inclusion morphisms to the co-product Knowledge.

Paper 2

Under development

Comparing Visualizations: Equivalence in
Perceptualization Processes

Paul Vickers, Nick Rossiter, Michael
Brockway, and Joe Faith.

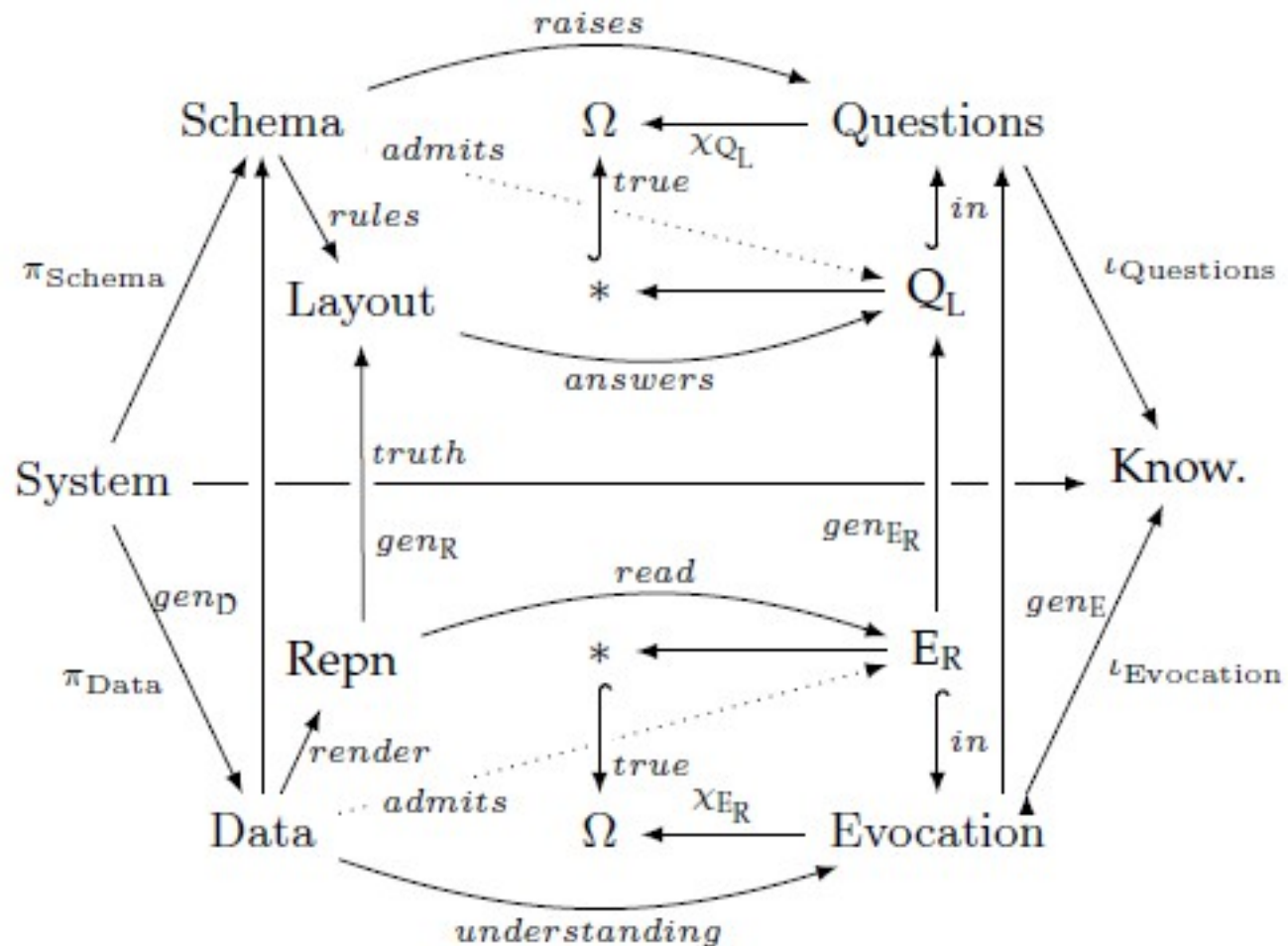


Fig. 5. The full visualization category with the subobject classifier furnishings added.

Basically a lattice of pullbacks

Formal Topos, Paper 2

Incomplete, Qs to be added

N name, T type,
V value, E encodings
P physical representation

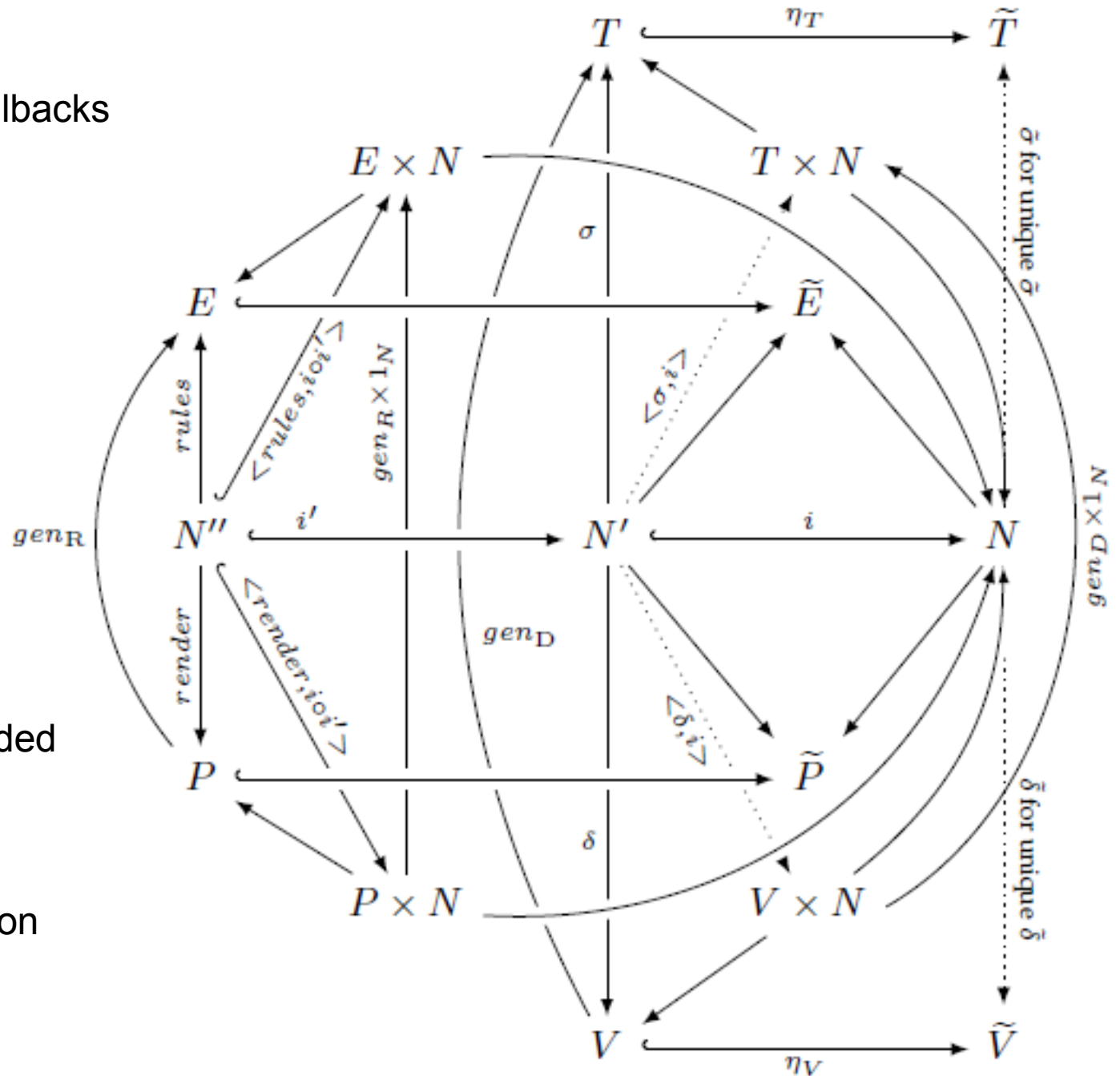


Fig. 7. Schema, Layout, Data, and Representation: The Expanded View

Final Thoughts

- Topos does look a very promising candidate for information systems
 - Ticks all the boxes for requirements
 - Well developed theoretically
- But some significant problems remain in application:
 - Need to develop an accessible way of drawing them
 - Need to develop how the internal logic will work
 - Need to make the internal logic accessible